

Economics of Networks

Introduction to Game Theory: Part 1

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Agenda

Decision theory:

- Theory of rational choice
- Choice under uncertainty

Games

- Definitions
- Strategies and best responses

Pure strategy Nash Equilibrium

Examples

Suggested reading: EK chapter 6; Osborne chapters 1-3

Motivation

Choice is a big part of economics

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How do we model this?

Rational Choice

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Three step process of rational choice:

- What is desirable?
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Observations:

- Note desires precede recognition of feasible alternatives
- Economics says nothing about what agent should desire

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A rational agent has a *preference ordering* over X

- For each $x, y \in X$, either $x \succeq y$ or $x \preceq y$
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Choice problem $C \subseteq X$

- Choose \succeq -maximal element of C

Utility Representation

Associate a number $u(x)$ to each $x \in X$, *utility* of x

- $u(x) \geq u(y)$ if and only if $x \succeq y$
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With uncertainty, we need *cardinal* information

- How do I compare outcomes with different probabilities?

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Examples:

- An even money gamble: with probability $\frac{1}{2}$ you lose \$10, with probability $\frac{1}{2}$ you win \$10
- Getting to work: with probability $\frac{9}{10}$ it takes you 20 minutes, with probability $\frac{1}{10}$ there is road work and it takes an hour

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- Independence: if $L \succ M$, then for any N and $p \in (0, 1]$, we have

$$pL + (1 - p)N \succ pM + (1 - p)N$$

Choice Under Uncertainty

Theorem

Suppose preferences \preceq satisfy the vNM axioms. There exists a utility function u on the set of outcomes Y such that $L \preceq M$ if and only if

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Expected utility theory

Suppose action a induces distribution $F^a(y)$ over consequences, expected utility

$$U(a) = \int u(y) dF^a(y)$$

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Notes:

- Objective versus Subjective uncertainty
- Savage (1954) axiomatizes subjective probability and subjective expected utility

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- Benefit from working hard on group work depends on others' efforts
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- How hard should I work?
- What route should I take?
- Should I tell a friend a secret?

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Do you work or shirk?

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What price do you set?

Normal Form Games

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Elements of a game

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- Set of actions or strategies
- Payoffs

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Player order, multiple moves, and information sets captured in extensive form games

- To come later

Normal Form Games

Definition (Normal Form Game)

A normal form game is a triple $(N, \{S_i\}_{i \in N}, \{u_i\}_{i \in N})$ such that

- $N = \{1, 2, \dots, n\}$ is a set of players
- S_i is the set of actions available to player i
- $u_i : S \rightarrow \mathbb{R}$ is the payoff of player i , where $S = \prod_{i \in N} S_i$ is the set of all action profiles

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Some notation:

- s_{-i} : vector of actions for all players except i
- S_{-i} : set of action profiles for all players except i
- $s = (s_i, s_{-i}) \in S$ is an action profile, or outcome

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Distinction also exists in normal form games when we talk about mixed strategies

The “Solution” of a Game

How do people play the game?

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How should people play the game?

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How do people play the game?

How should people play the game?

Large number of “solution concepts” based on different assumptions about

- What players know about each others' plans
- How smart players are
- How smart players think others are

Dominant Strategies

Example: The Prisoner's Dilemma

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Confess	$(-3, -3)$	$(0, -4)$
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“Confess” **dominates** “silence”

Dominant Strategy Equilibrium

A fairly compelling solution concept: everyone plays a dominant strategy

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Definition

A strategy $s_i^* \in S_i$ is *dominant* for player i if for all $s_i \in S_i$ and all $s_{-i} \in S_{-i}$

$$u_i(s_i^*, s_{-i}) \geq u_i(s_i, s_{-i})$$

A strategy profile s^* is a *dominant strategy equilibrium* if s_i^* is a dominant strategy for each i .

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Issue: this rarely exists

Dominated Strategies

Conversely, we might think to eliminate strategies that are dominated

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A strategy $s_i \in S_i$ is *strictly dominated* if there exists $s'_i \in S_i$ such that for all $s_{-i} \in S_{-i}$

$$u_i(s'_i, s_{-i}) > u_i(s_i, s_{-i})$$

A strategy $s_i \in S_i$ is *weakly dominated* if there exists $s'_i \in S_i$ such that for all $s_{-i} \in S_{-i}$

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Iterated Deletion

No one should play a dominated strategy

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No dominant strategy

- Still dominance solvable

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More formally, define the iterative procedure:

- Step 0: Define $S_i^0 = S_i$ for each i
- Step $k > 0$: Define for each i

$$S_i^k = \left\{ s_i \in S_i^{k-1} \mid \nexists s'_i \in S_i^{k-1} \text{ that dominates } s_i \right\}$$

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The set S^∞ of strategy profiles is what survives

Iterated Deletion

Theorem

Suppose that either:

- *S_i is finite for each i , or*
- *$u_i(s_i, s_{-i})$ is continuous and S_i is compact for each i .*

Then S_i^∞ is nonempty for each i .

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May not yield a unique prediction

Best Responses

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Should play a best reply to *some* conjecture

- What conjectures should a player entertain?

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Why might this be reasonable?

Interpretation of Nash Equilibrium

Two main justifications:

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Nash Equilibrium is a standard workhorse in economic models

- Might not be reasonable in all contexts

Example: Bertrand Competition

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- Marginal cost $c > 0$
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- Yes! Both firms earn zero profit, no way to improve.

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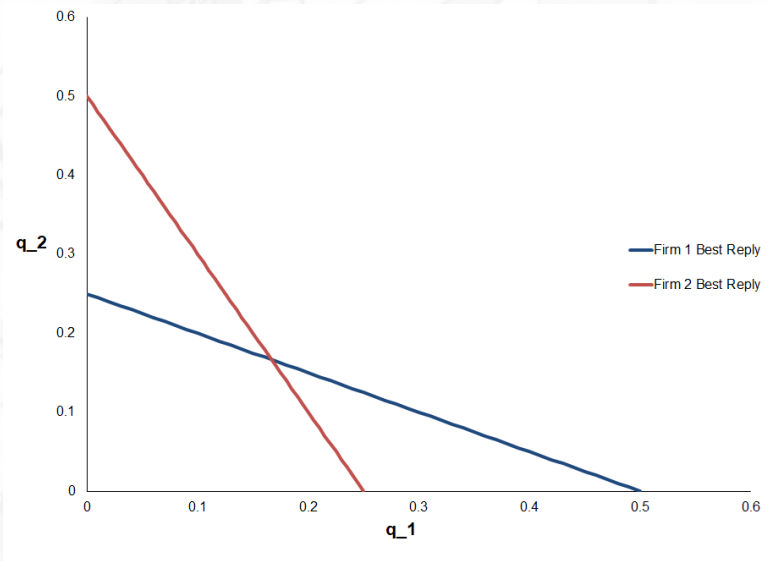
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Unique Nash Equilibrium:

$$q_1^* = q_2^* = \frac{2(1 - c)}{3}$$

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Example: The Partnership Game

Recall our earlier choice to work or shirk:

	Work	Shirk
Work	(2, 2)	(-1, 1)
Shirk	(1, -1)	(0, 0)

No dominant or dominated strategies

Best reply to work is work, best reply to shirk is shirk

- Two pure strategy Nash Equilibria
- Outcome depends on conjectures

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Refinement is hard

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	Heads	Tails
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No pure strategy Nash Equilibrium exists

- Next time: mixed strategy equilibrium